New Directions in Automated Mechanism Design

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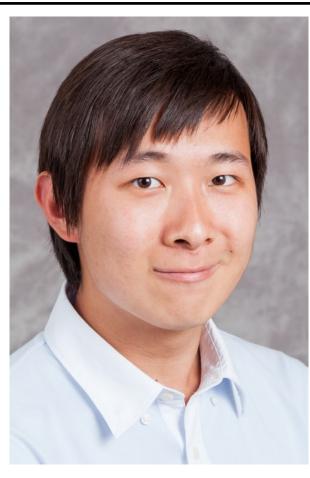


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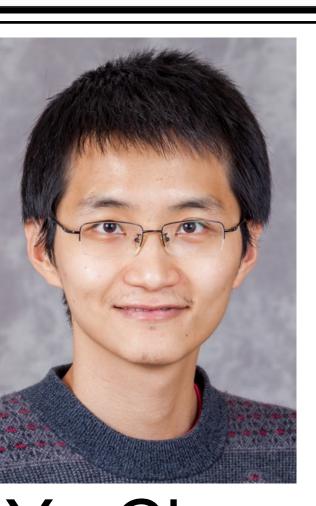


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Mechanism design

Make decisions based on the preferences (or other information) of one or more agents (as in social choice)

Focus on *strategic* (game-theoretic) agents with *privately held* information; have to be *incentivized* to reveal it truthfully

Popular approach in design of auctions, matching mechanisms, ...

Sealed-bid auctions (on a single item)



Bidder i determines how much the item is worth to her (v_i)

Writes a bid (v_i) on a piece of paper

How would you bid? How much would I make?

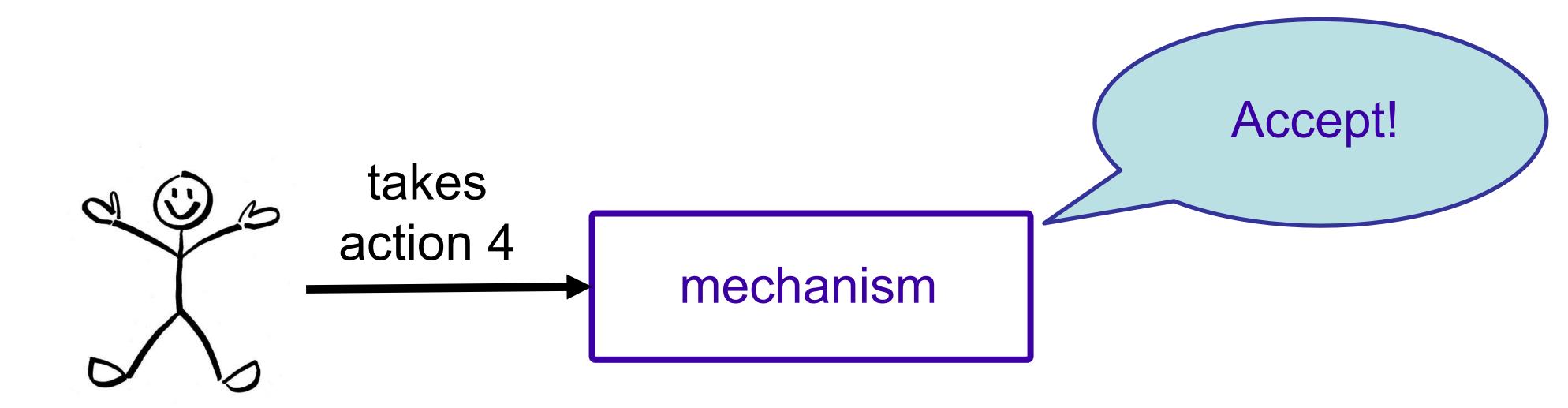
First price: Highest bid wins, pays bid

Second price: Highest bid wins, pays next-highest bid

First price with reserve: Highest bid wins iff it exceeds *r*, pays bid **Second price with reserve:** Highest bid wins iff it exceeds *r*, pays next highest bid or *r* (whichever is higher)

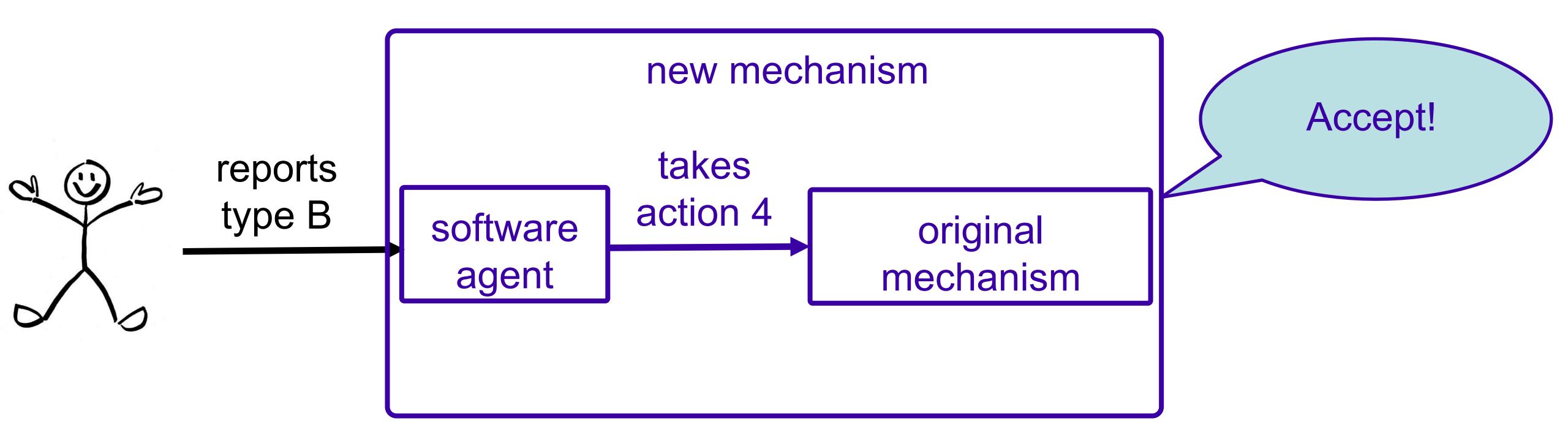
Revelation Principle

Anything you can achieve, you can also achieve with a truthful (AKA incentive compatible) mechanism.



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Anything you can achieve, you can also achieve with a truthful (AKA incentive compatible) mechanism.



Automated mechanism design input

Instance is given by

Set of possible outcomes

Set of agents

For each agent set of possible *types*

probability distribution over these types

Objective function

Gives a value for each outcome for each combination of agents' types E.g., social welfare, revenue

Restrictions on the mechanism

Are payments allowed?

Is randomization over outcomes allowed?

What versions of incentive compatibility (IC) & individual rationality (IR) are used?

How hard is designing an optimal deterministic mechanism (without reporting costs)?

[C. & Sandholm UAI'02, ICEC'03, EC'04]

NP-complete (even with 1 reporting agent):	Solvable in polynomial time (for any constant number of agents):
1.Maximizing social welfare (no payments)	1. Maximizing social welfare (not regarding the
2. Designer's own utility over outcomes (no payments)	payments) (VCG)
3. General (linear) objective that doesn't regard payments	
4.Expected revenue	

1 and 3 hold even with no IR constraints

Positive results (randomized mechanisms)

[C. & Sandholm UAI'02, ICEC'03, EC'04]

- Use linear programming
- Variables:

```
p(o \mid \theta_1, ..., \theta_n) = probability that outcome o is chosen given types <math>\theta_1, ..., \theta_n (maybe) \pi_i(\theta_1, ..., \theta_n) = i's payment given types \theta_1, ..., \theta_n
```

• Strategy-proofness constraints: for all i, θ_1 , ... θ_n , θ_i :

$$\Sigma_{o}p(o \mid \theta_{1}, ..., \theta_{n})u_{i}(\theta_{i}, o) + \pi_{i}(\theta_{1}, ..., \theta_{n}) \geq \Sigma_{o}p(o \mid \theta_{1}, ..., \theta_{i}', ..., \theta_{n})u_{i}(\theta_{i}, o) + \pi_{i}(\theta_{1}, ..., \theta_{i}', ..., \theta_{n})$$

• Individual-rationality constraints: for all i, θ_1 , ... θ_n :

$$\Sigma_{o}p(o \mid \theta_{1}, ..., \theta_{n})u_{i}(\theta_{i}, o) + \pi_{i}(\theta_{1}, ..., \theta_{n}) \geq 0$$

Objective (e.g., sum of utilities)

$$\Sigma_{\theta_1, ..., \theta_n} p(\theta_1, ..., \theta_n) \Sigma_i (\Sigma_o p(o | \theta_1, ..., \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, ..., \theta_n))$$

- Also works for BNE incentive compatibility, ex-interim individual rationality notions, other objectives, etc.
- For deterministic mechanisms, can still use mixed integer programming: require probabilities in {0, 1}
 - -Remember typically designing the optimal deterministic mechanism is NP-hard

A simple example

One item for sale (free disposal)

2 agents, IID valuations: uniform over {1, 2}

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

(What is optimal expected welfare?)

Agent 1's valuation

0.25 0.25 0.25 0.25

probabilities

Agent 2's valuation

Status: OPTIMAL

Objective: obj = 1.5 (MAXimum)

[nonzero variables:]

p t 1 1 o3

(probability of disposal for (1, 1))

(probability 1 gets the item for (2, 1))

(probability 2 gets the item for (1, 2))

t 2 2 o2

(probability 2 gets the item for (2, 2)) (1's payment for (2, 2))

pi 2 2 1

pi 2 2 2

4

(2's payment for (2, 2))

Our old AMD solver [C. & Sandholm, 2002, 2003] gives:

A slightly different distribution

One item for sale (free disposal)

2 agents, valuations drawn as on right

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

(What is optimal expected welfare?)

Agent 1's valuation

Agent 2's valuation		
1		
0.251	0.250	
0.00	0 0 40	

probabilities

Status: OPTIMAL

Objective: obj = 1.749 (MAXimum)

[some of the nonzero payment variables:]

You'd better be really sure about your distribution!

A nearby distribution without correlation

One item for sale (free disposal)

2 agents, valuations IID: 1 w/ .501, 2 w/ .499

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

Agent 1's valuation

(What is optimal expected welfare?)

Agent 2's valuation

1 2

0.251001 0.249999

0.249999 0.249001

probabilities

Status: OPTIMAL

Objective: obj = 1.499 (MAXimum)

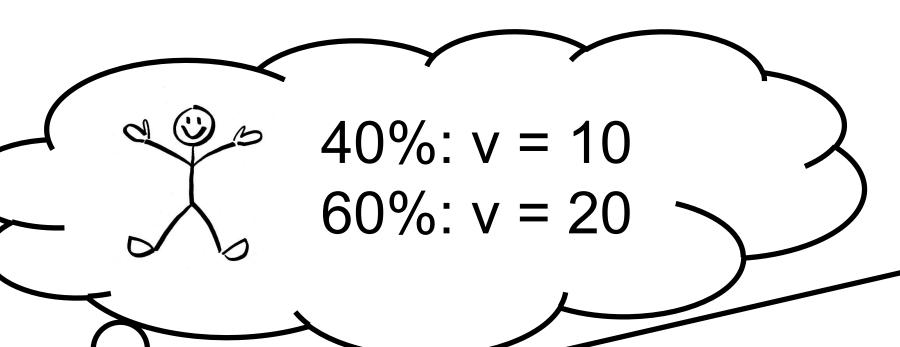
Cremer-McLean [1985]

For every agent, consider the following matrix Γ of conditional probabilities, where Θ is the set of types for the agent and Ω is the set of signals (joint types for other agents, or something else observable to the auctioneer)

$$\Gamma = \begin{bmatrix} \pi(1|1) & \cdots & \pi(|\Omega||1) \\ \vdots & \ddots & \vdots \\ \pi(1||\Theta|) & \cdots & \pi(|\Omega|||\Theta|) \end{bmatrix}$$

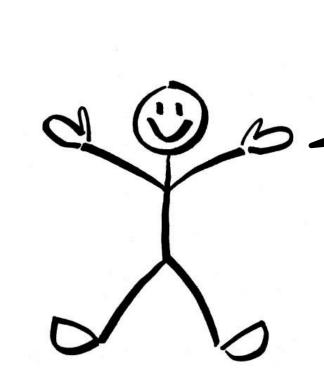
If Γ has rank $|\Theta|$ for every agent then the auctioneer can allocate efficiently and extract the full surplus as revenue (!!)

Standard setup in mechanism design



(1) Designer has beliefs about agent's *type* (e.g., preferences)

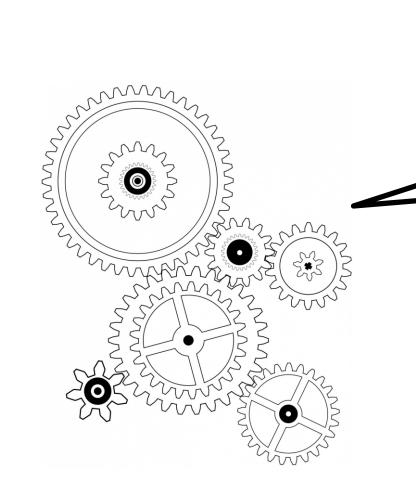
(2) Designer announces mechanism (typically mapping from reported types to outcomes)



(3) Agent *strategically* acts in mechanism (typically type report), *however she*

likes at no cost

y = 20



(4) Mechanism functions as specified

The mechanism may have more information about the specific agent!

application

online marketplaces selling insurance university admissions webpage ranking

information

actions taken online driving record courses taken links to page

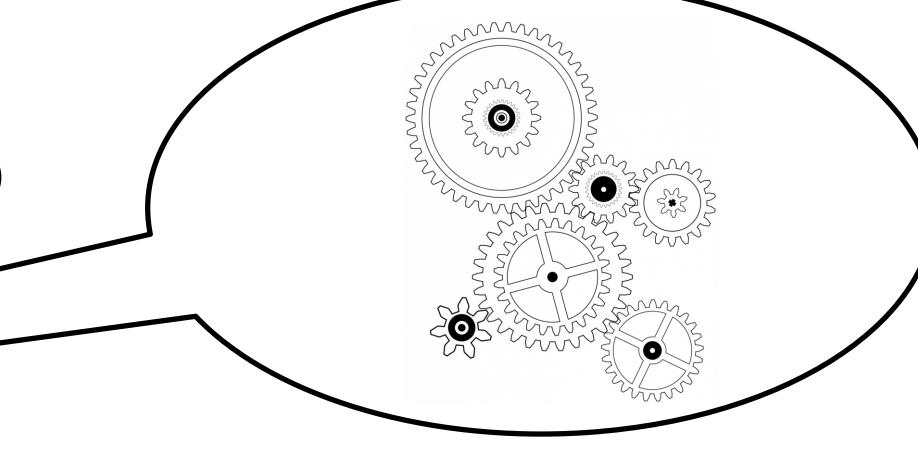


Attempt 1 at fixing this



30%: v = 10

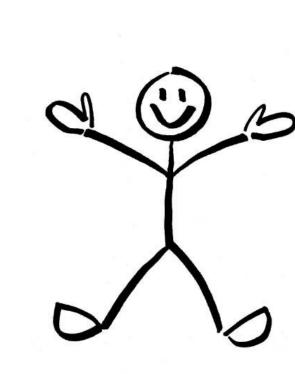
70%: v = 20



(0) Agent acts in the world (naively?)

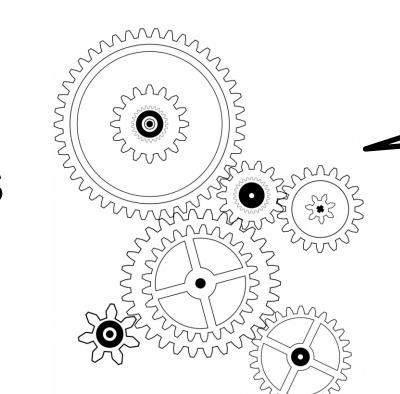
(1) Designer obtains beliefs about agent's *type* (e.g., preferences)

(2) Designer announces mechanism (typically mapping from reported types to outcomes)



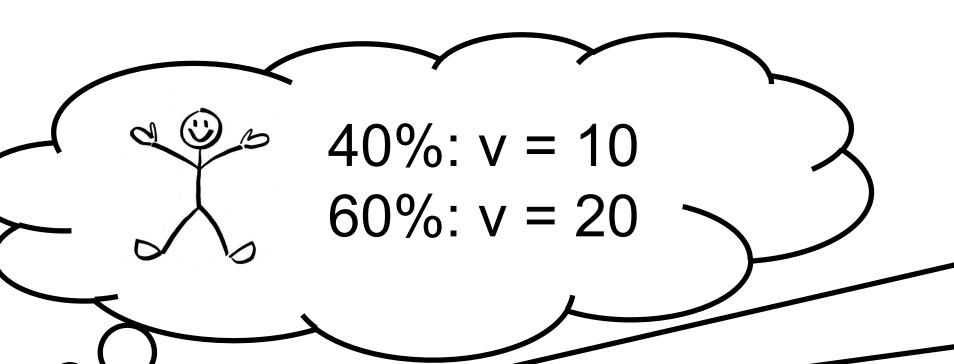
y = 20

(3) Agent strategically acts in mechanism (typically type report), however she likes at no cost



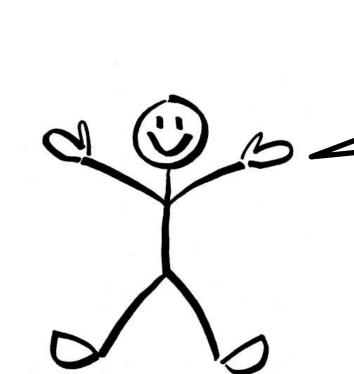
(4) Mechanism functions as specified

Attempt 2: Sophisticated agent



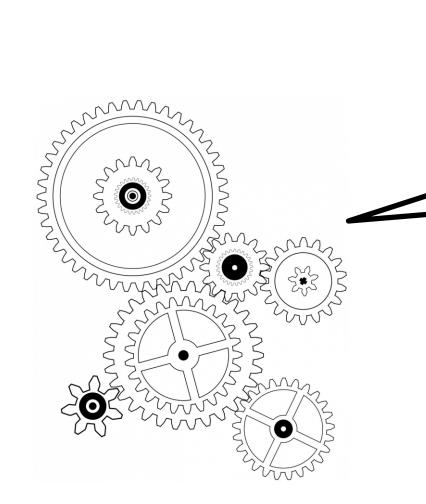
(1) Designer has prior beliefs about agent's *type* (e.g., preferences)

(2) Designer announces mechanism (typically mapping from reported types to outcomes)



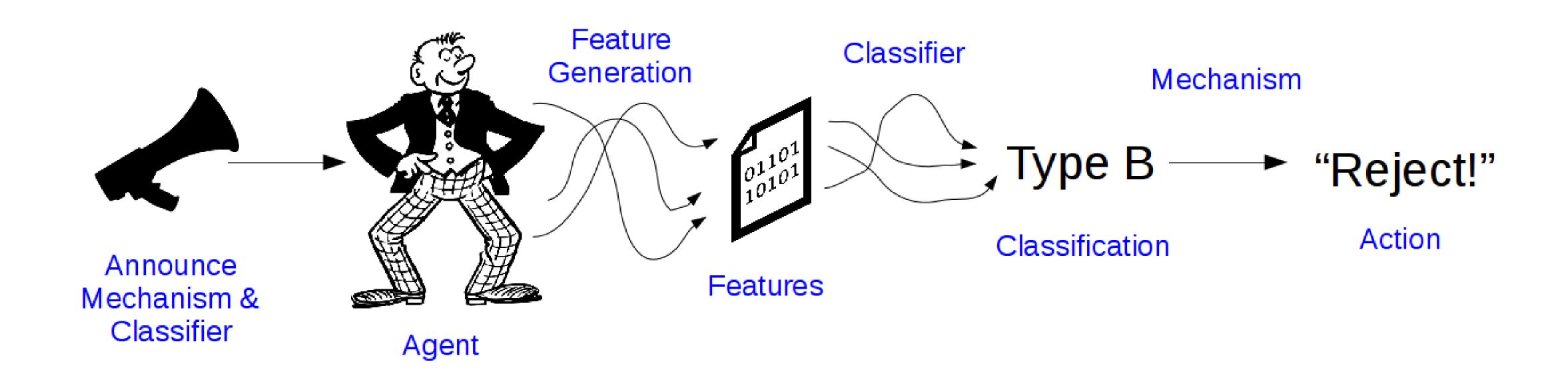
Show me pictures of cats v = 20

(3) Agent strategically takes possibly costly actions



(4) Mechanism functions as specified

Machine learning view



See also later work by Hardt, Megiddo, Papadimitriou, Wootters [2015/2016]

From Ancient Times...

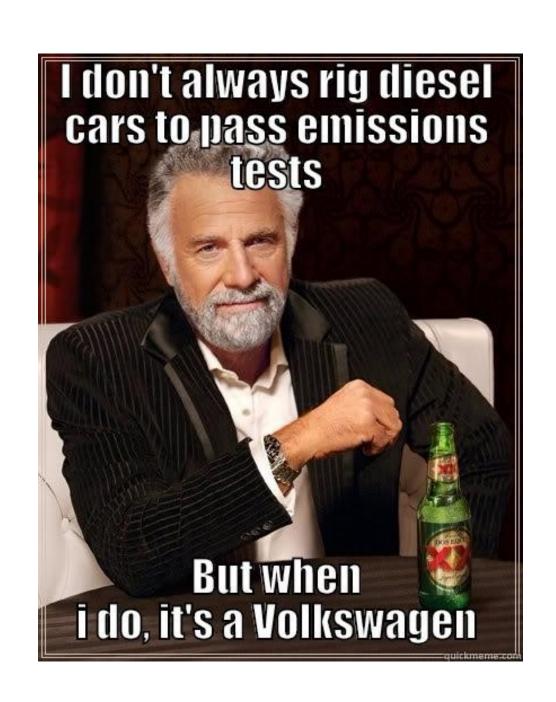


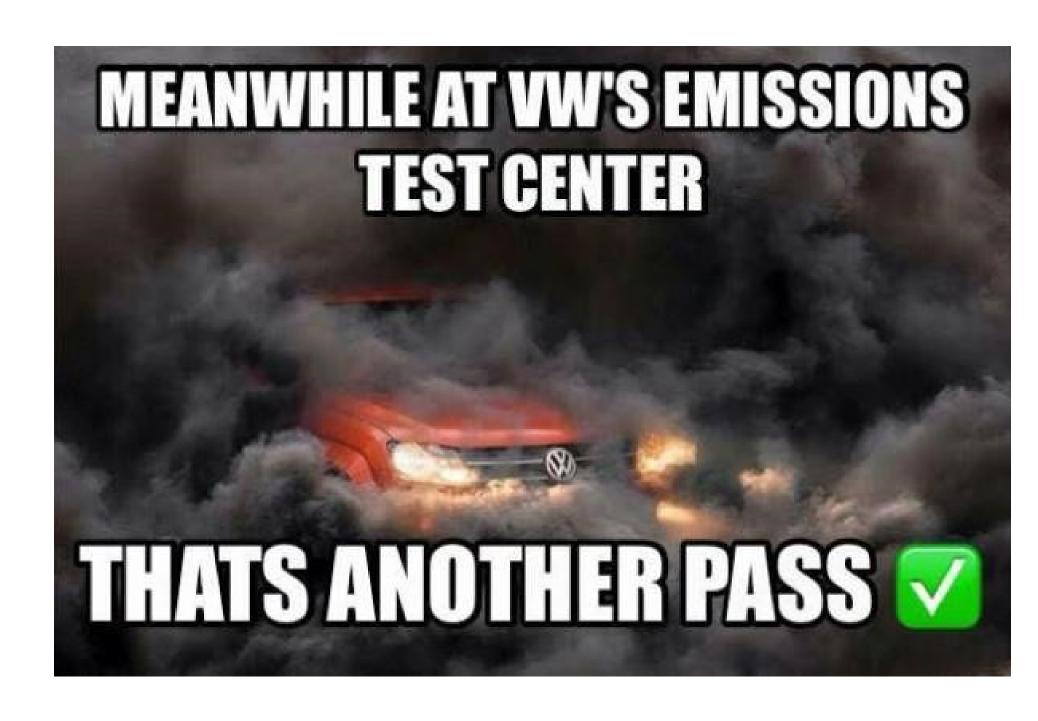


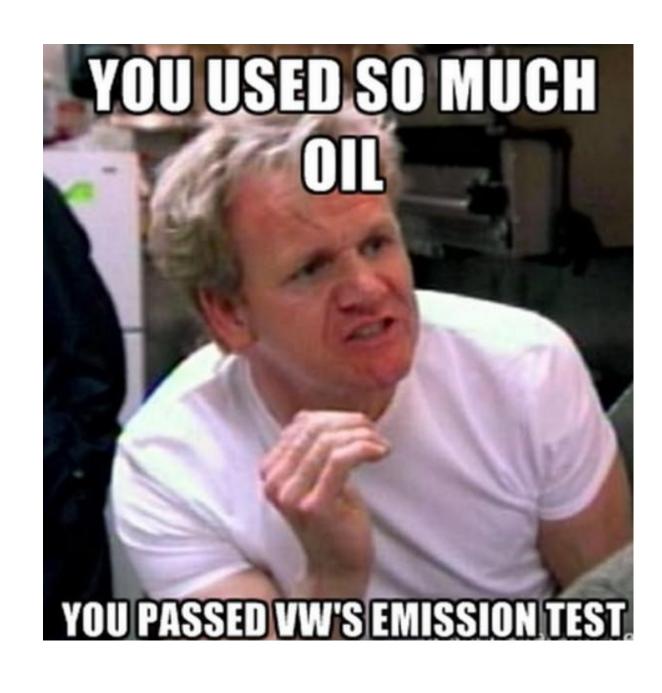
Jacob and Esau

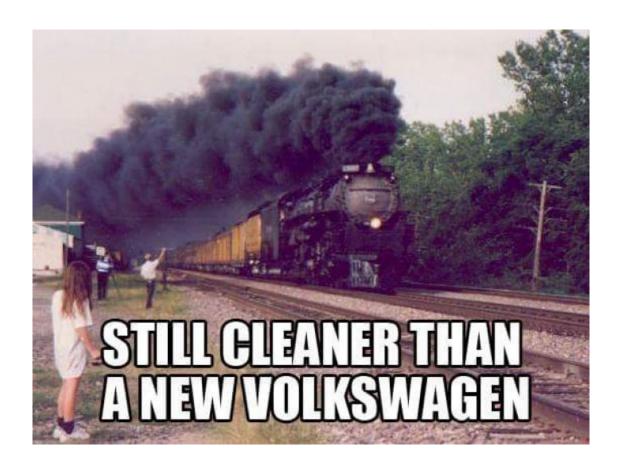
Trojan Horse

... to Modern Times









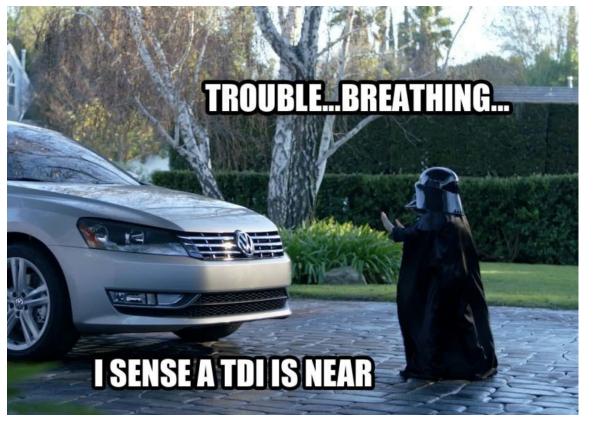






Illustration: Barbara Buying Fish From Fred

Classifications ($\hat{t} \in \hat{T}$): \hat{fresh} , \hat{ok} , $rot\hat{t}en$





... continued

Effort Function $(E: T \times \hat{T} \to \mathbb{R})$:

 $fresh \\ ok \\ rotten$

\hat{fresh}	\hat{ok}	$rot \hat{t}en$
0	0	0
10	0	0
30	10	0



Valuation Function $(V: T \times A \rightarrow \mathbb{R})$:

$$V(\cdot, accept) = 20, \ V(\cdot, reject) = 0$$

$\underline{\mathbf{Mechanism}\ M: \hat{T} \to A}$

First Try: $M = \hat{fresh} \rightarrow accept, \ \hat{ok} \rightarrow accept, \ rot\hat{ten} \rightarrow reject$

... continued

Effort Function $(E: T \times \hat{T} \to \mathbb{R})$:

fresh ok rotten

\hat{fresh}	\hat{ok}	$rot \hat{t}en$
0	0	0
10	0	0
30	10	0



Valuation Function $(V: T \times A \rightarrow \mathbb{R})$:

 $V(\cdot, accept) = 20, \ V(\cdot, reject) = 0$

$\underline{\mathbf{Mechanism}\ M: \hat{T} \to A}$

First Try: $M = \hat{fresh} \rightarrow accept, \ \hat{ok} \rightarrow accept, \ rot\hat{ten} \rightarrow reject$

Better: $M^* = \hat{rresh} \rightarrow accept$, $\hat{ok} \rightarrow reject$, $rot\hat{t}en \rightarrow reject$.

Comparison With Other Models

 $fresh\\ok\\rotten$

\hat{fresh}	\hat{ok}	$rot \hat{t}en$
0	0	0
0	0	0
0	0	0

Standard Mechanism Design

fresh ok rotten

fresh	ok	rotten
0	∞	0
∞	0	0
0	0	0

Mechanism Design with Partial Verification

fresh ok rotten

\hat{fresh}	\hat{ok}	$rot \hat{t}en$
∞	0	0
1.2	5	-∞
-3	0	0

Mechanism Design with Signaling Costs

Green and Laffont. Partially verifiable information and mechanism design. RES 1986 Auletta, Penna, Persiano, Ventre. Alternatives to truthfulness are hard to recognize. AAMAS 2011

Question

Given:

Types $(t \in T)$ Actions $(a \in A)$ Choice Function $(F : T \to A)$

()	((
fresh	accept	$fresh \rightarrow accept$
ok	reject	$ok \rightarrow accept$
rotten		$rotten \rightarrow reject$

Classifications ($\hat{t} \in \hat{T}$): \hat{fresh} , \hat{ok} , $rot\hat{t}en$





Effort Function($E: T \times \hat{T} \to \mathbb{R}$):

	\hat{fresh}	\hat{ok}	$rot \hat{t}en$
fresh	0	0	0
ok	10	0	0
rotten	30	10	0



Valuation Function $(V: T \times A \rightarrow \mathbb{R})$:

$$V(\cdot, accept) = 20, \ V(\cdot, reject) = 0$$

Then:

Does there exist a

Mechanism $M: \hat{T} \to A$

which implements the choice function?

NP-complete!

Auletta, Penna, Persiano, Ventre. Alternatives to truthfulness are hard to recognize. AAMAS 2011

Results



with Andrew Kephart (AAMAS 2015)

		Transfers (T)		No Transfers (NT)	
		Two Outcomes (TO)	Injective SCF (FI)	Two Outcomes (TO)	Injective SCF (FI)
Free Utilities (FU)	Unrestricted Costs (U) $\{0, \infty\}$ Costs (ZI)	NP-c NP-c	NP-c NP-c	NP-c NP-c	NP-c P
Targeted Utilities (TU)	Unrestricted Costs (U) $\{0, \infty\}$ Costs (ZI)	NP-c NP-c	P P	NP-c NP-c	P P

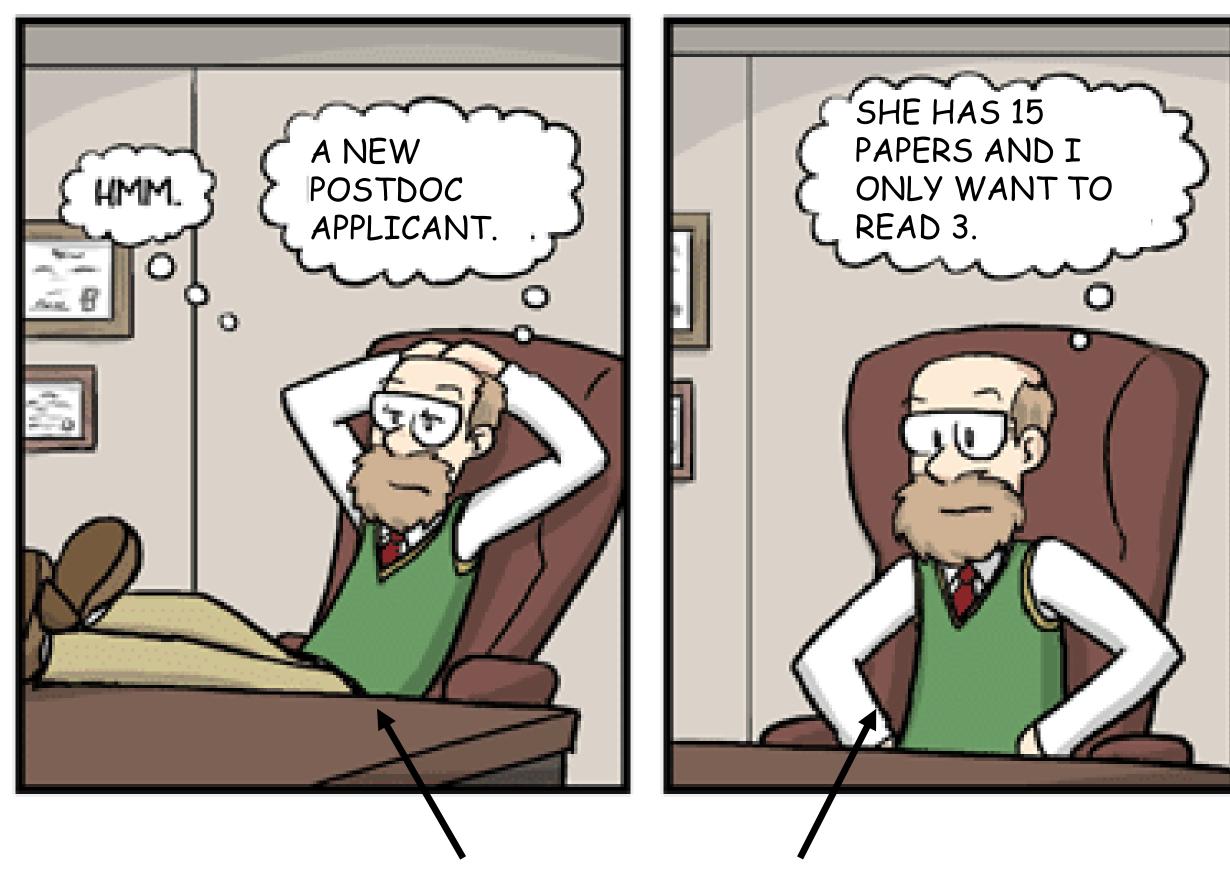
Non-bolded results are from:

Auletta, Penna, Persiano, Ventre. Alternatives to truthfulness are hard to recognize. AAMAS 2011

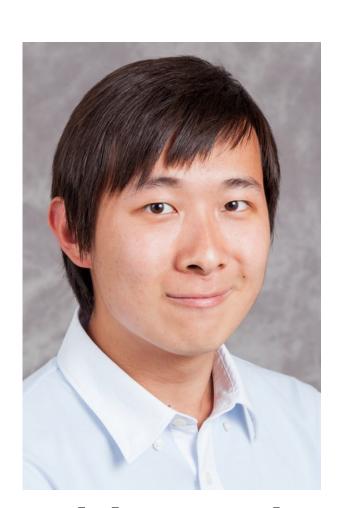
Hardness results fundamentally rely on revelation principle failing – conditions under which revelation principle still holds in Green & Laffont '86 and Yu '11 (partial verification), and Kephart & C. EC'16 (costly signaling).

When Samples Are Strategically Selected

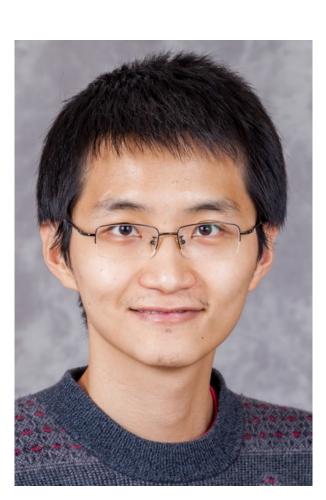
ICML 2019, with







Hanrui
Zhang
(Duke)



Yu Cheng (Duke → IAS → UIC)

Academic hiring...



Charlie, Bob's student

Academic hiring...

I NEED TO CHOOSE THE BEST 3 PAPERS TO CONVINCE BOB, SO THAT HE WILL HIRE ALICE.

CHARLIE WILL DEFINITELY PICK THE BEST 3 PAPERS BY ALICE, AND I NEED TO CALIBRATE FOR THAT.



A distribution (Alice) over paper qualities $\theta \in \{g, b\}$ arrives, which can be either a good one $(\theta = g)$ or a bad one $(\theta = b)$



Alice, the postdoc applicant

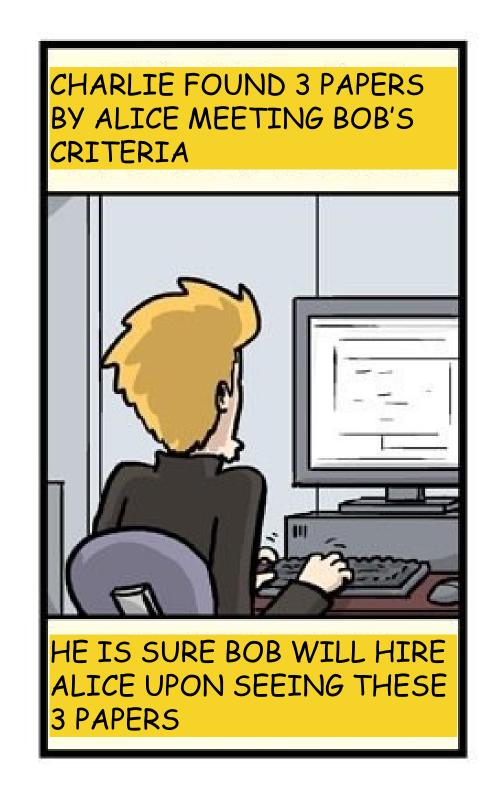
The principal (Bob) announces a policy, according to which he decides, based on the report of the agent (Charlie), whether to accept θ (hire Alice)



The agent (Charlie) has access to n(=15) iid samples (papers) from θ (Alice), from which he chooses m(=3) as his report



The agent (Charlie) sends his report to the principal, aiming to convince the principal (Bob) to accept θ (Alice)



The **principal (Bob)** observes the **report** of the **agent** (**Charlie**), and makes the decision according to the policy announced

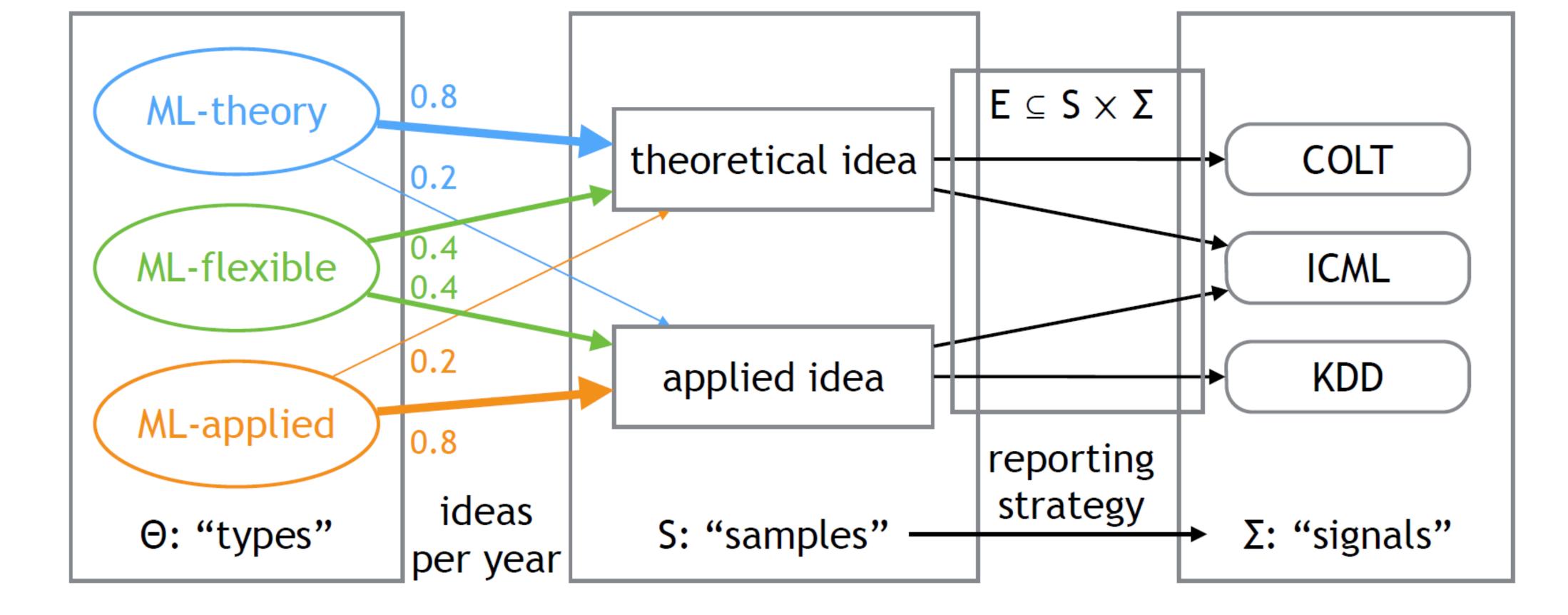






Questions

How does strategic selection affect the principal's policy? Is it easier or harder to classify based on **strategic samples**, compared to when the principal has access to **iid samples**? Should the principal ever have a **diversity** requirement (e.g., at least 1 mathematical paper and at least 1 experimental paper), or only go by total quality according to a single metric?



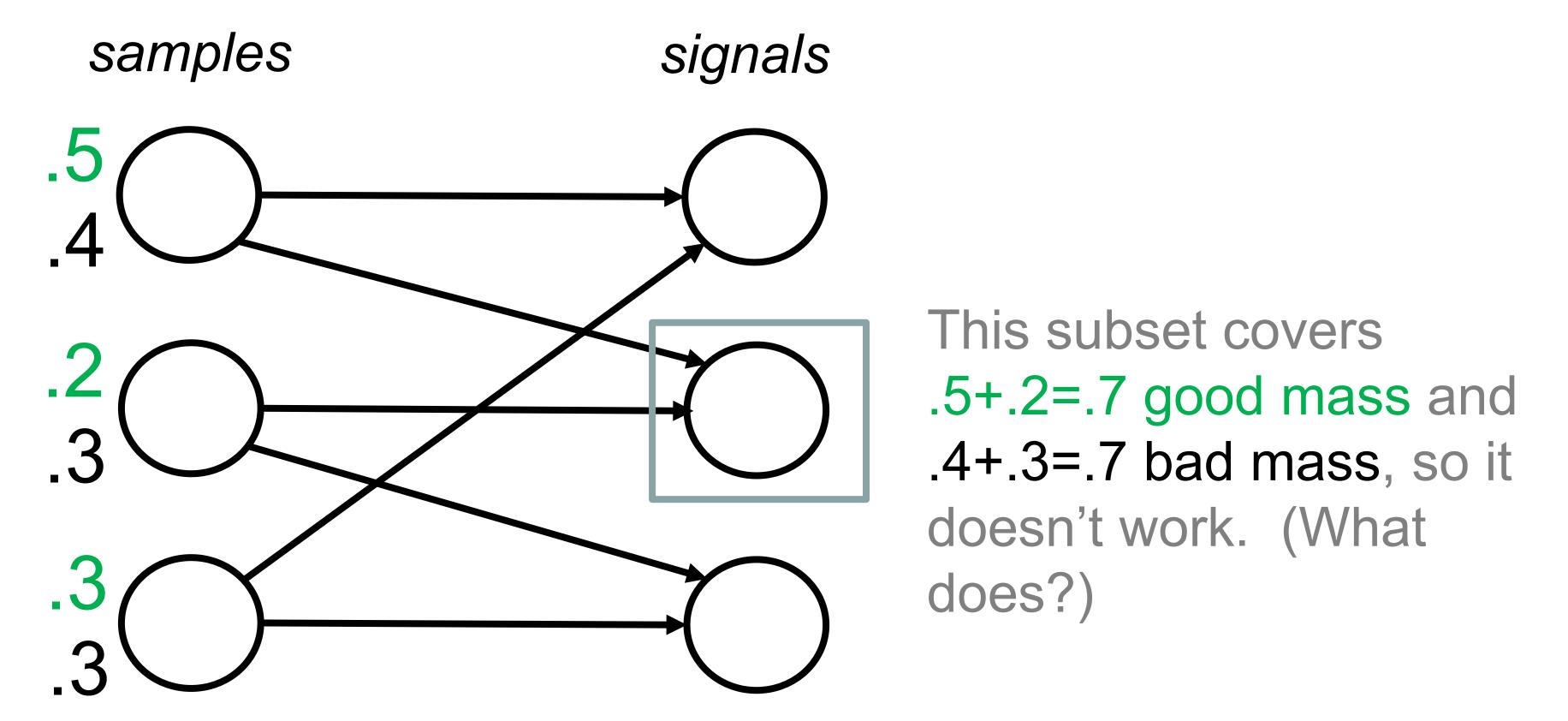
Agent's problem:

- "How do I distinguish myself from other types?"
- "How many samples do I need for that?"

Principal's problem:

- "How do I tell ML-flexible agents from others?"
- "At what point in their career can I reliably do that?"

One good and one bad distribution

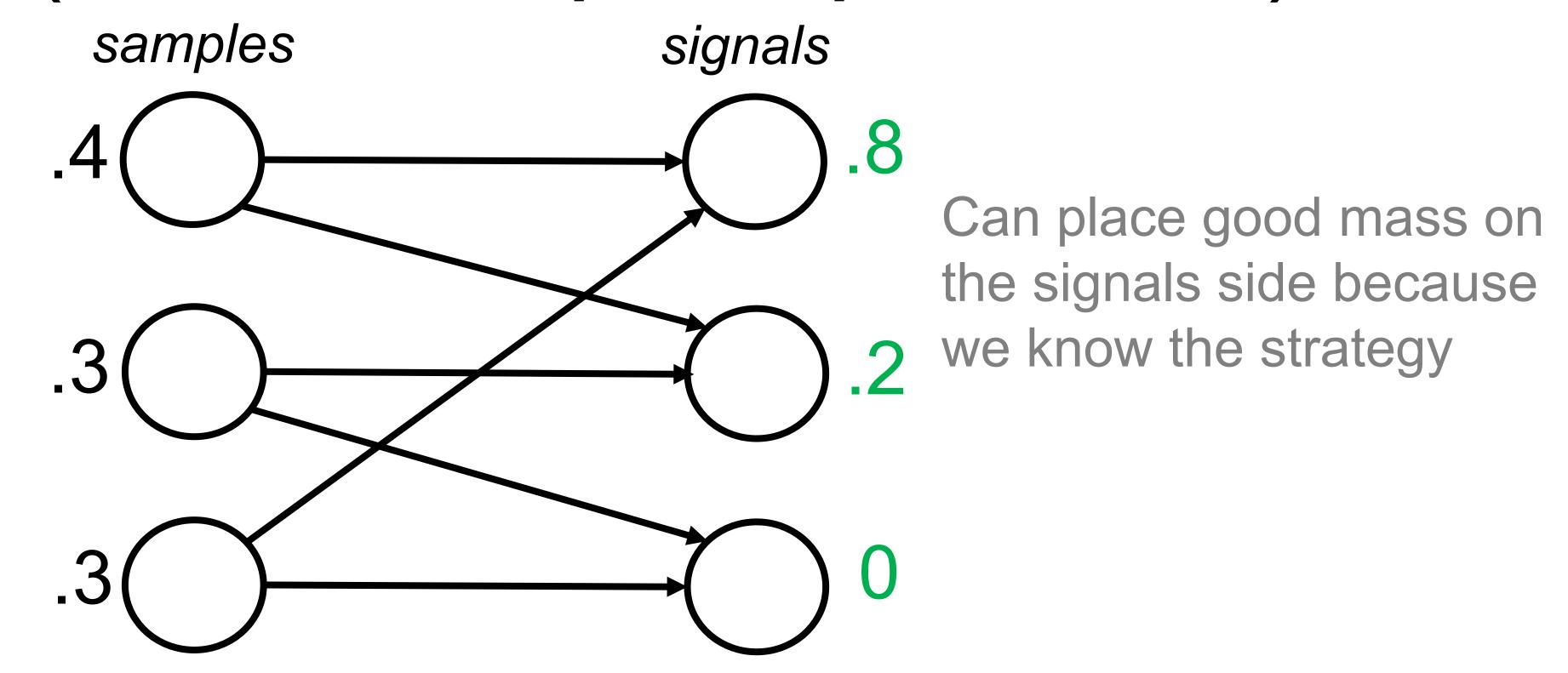


Pick a subset of the right-hand side (to accept) that maximizes (green mass covered - black mass covered)

If positive, can (eventually) distinguish; otherwise not.

NP-hard in general.

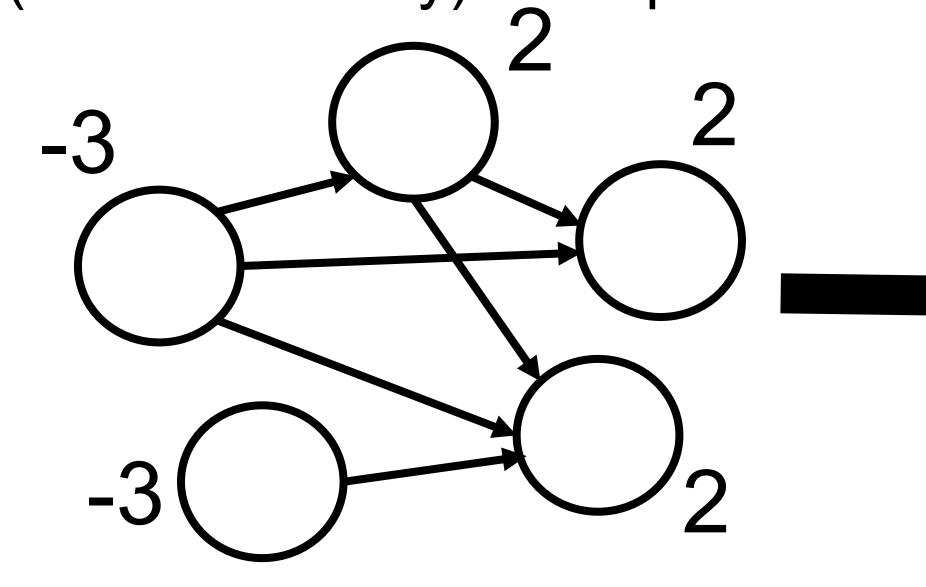
But if we know the strategy for the good distribution (revelation principle holds):



Solve as maximum flow/matching from left to right with capacities on vertices Duality gives set of signals to accept (~Hall's marriage theorem)

Optimization: reduction to min cut (when revelation principle holds)

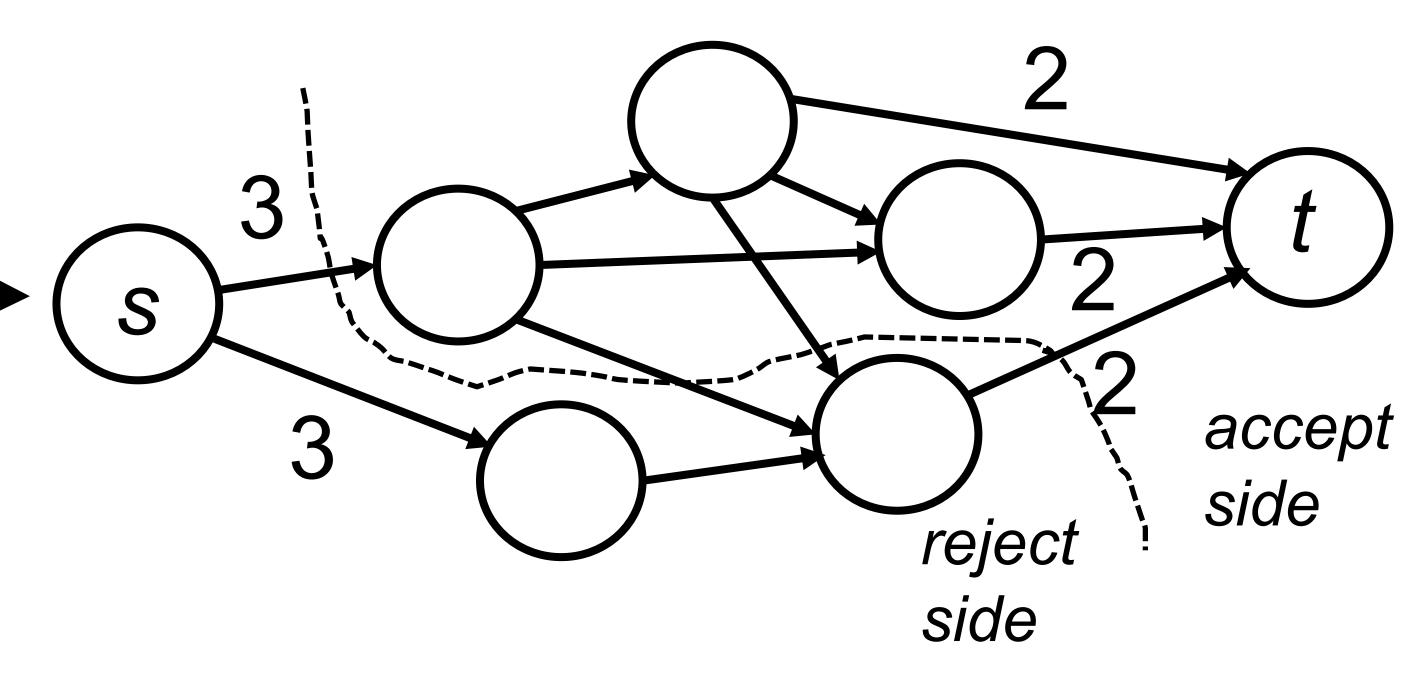
types are vertices; edges imply ability to (cost-effectively) misreport



In sampling case, can check existence of edges with previous technique

Values are P(type)*value(type)

edges between types have capacity ∞



Can be generalized to more outcomes than accept/reject, if types have the same utility over them.

Conclusion

First part:

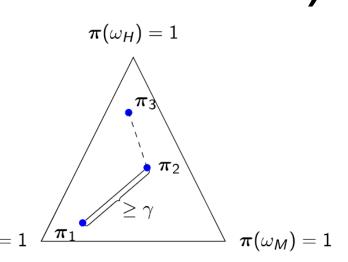
When considering correlation, small changes can have a huge effect

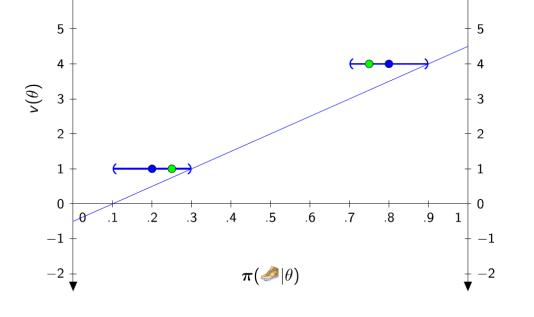
Automatically designing robust mechanisms addresses this

Combines well with learning (under some conditions)

0.251	0.250
0.250	0.249

	0.251001	0.249999	
V .	0.249999	0.249001	





Second part:

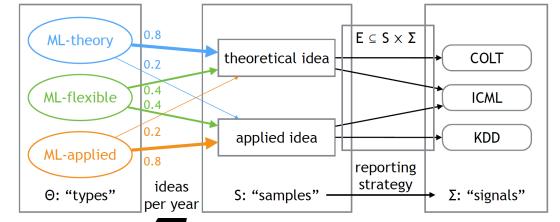
With costly or limited misreporting, revelation principle can fail

Causes computational hardness in general

Sometimes agents report based on their samples

Some efficient algorithms for the infinite limit case; sample bounds

Effort Function $(E: T \times \hat{T} \to \mathbb{R})$:				
	\hat{fresh}	\hat{ok}	$rot \hat{t}en$	
$fresh \\ ok$	О	0	0	
	10	0	0	
rotten	30	10	0	



Thank you for your attention!